

# Brief Announcement: Enhancing Permissiveness of Transactional Memory via Time-Warp

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# Introduction

To guarantee a given correctness level, a TM aborts transactions.

```
function commit(Transaction tx):
```

```
  for each  $\langle datum, version \rangle \in tx.readSet$  do
```

```
    if not latestVersion(datum, version) then
```

```
      abort(tx)
```

# Problem at hands

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Abort  $T$  if its reads are not up-to-date when it attempts to commit.

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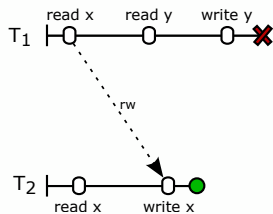
Abort  $T$  if its reads are not up-to-date when it attempts to commit.

Serializability:

- Necessary condition
- But **not** sufficient
- Deemed to be practical

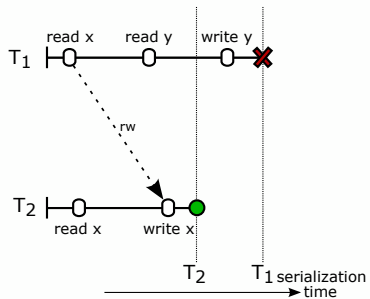
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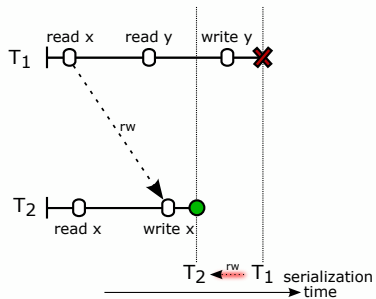




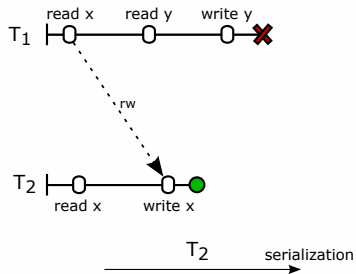
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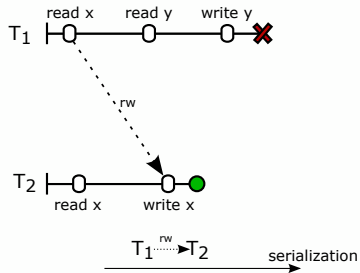
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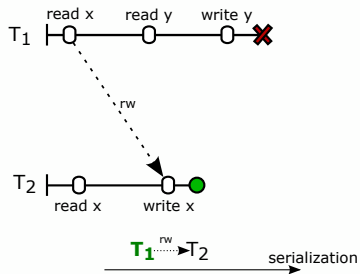
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**Time-warp** commit

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Notion of avoiding unnecessary aborts:

- If it only aborts a transaction when the resulting history (without the abort) does not respect some target correctness criterion
- State of the art TMs are far from being permissive
- Cost is **not** negligible.

# Goal

Enhance permissiveness without such costs

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Enhance permissiveness without such costs:

- More restrictive abort condition
- Always read consistently
- Wait-free read-only transactions (mv-permissiveness)

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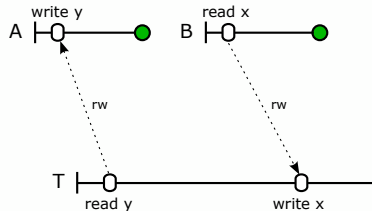
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Look out for a specific structure:

- Three transactions connected
  - ▶ a triad

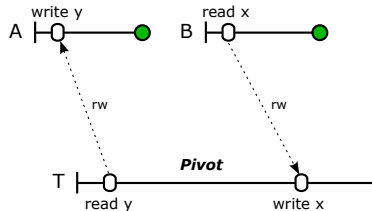


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Look out for a specific structure:

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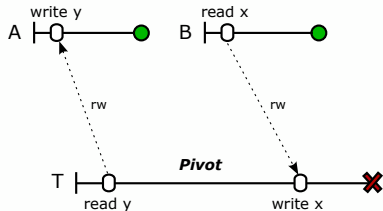


# Time-warp

When can we **not** apply this idea?

Look out for a specific structure:

- Three transactions connected
  - ▶ a triad
- The link between all three
  - ▶ the pivot
- Abort if:
  - ▶ Completes a triad
  - ▶ Whose pivot time-warp commits





# Theoretical and Practical results

- Virtual World Consistency

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  - ▶ Serializability for committed transactions
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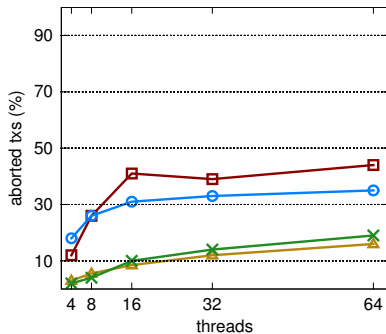
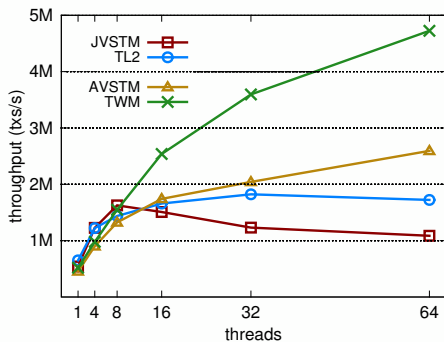
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  - ▶ Serializability for committed transactions
  - ▶ Prevent a range of phenomena for running transactions
- Lock-free prototype in Java: TWM
- 64 core machine
- Comparison with TL2, JVSTM, AVSTM

# Ordered Skip-List



Questions?

Thank you

Enhancing Permissiveness in Transactional Memory via Time-Warping,  
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