

Deterministic execution of TM applications

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Introduction

- Multicore systems and multithreaded applications
- Threads interleave in arbitrary order (nondeterminism)
- Difficult to develop, test, debug
- Difficult to provide fault-tolerance and attack-tolerance
- Determinism:
 - Repeatability
 - Easier finding and solving bugs
 - Easier providing tolerance to faults/attacks in replica-based systems with no communication among replicas

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Deterministic execution

- 1 running thread at a time
- Threads execute in a previously defined order (round-robin)
- Synchronization operations points where a thread changes its state
- Thread's states: running, ready, blocked
- Providing strong determinism deterministic execution of code within and outside critical sections/atomic blocks (important for code with data races)
- Implementations:
 - Det-Serial
 - Det-Parallel



Implementation – Det-Serial

- A runtime C library
- Synchronization operations points where a thread changes its state:
 - pthread_create, pthread_yield, pthread_join, pthread_barrier_init, pthread_barrier_destroy, pthread_barrier_wait

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- sleep

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• Running all code in serial

Implementation – Det-Parallel

- Based on Det-Serial
- Running non-transactional code in serial
- Transactions:
 - start in serial
 - run in parallel
 - commit in serial
- STM with 2 conflict detection policies (Lazy, Eager)
- Det-Parallel:
 - Lazy Det-Parallel
 - Eager Det-Parallel



Implementation – Det-Parallel

Lazy Det-Parallel Transactions: start in serial run in parallel commit in serial (conflict detection and aborts)

Eager Det-Parallel

- Transactions:
 - start in serial

Additional barriers to synchronize threads

- run in parallel (conflict detection and aborts)
- commit in serial

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Implementation – Det-Parallel(2)

Eager Det-Parallel

- Transactions:
 - start in serial
 - run in parallel (conflict detection and aborts)
 - commit in serial

Step 1.

- the first-to-commit transaction never aborts
- other conflicting transactions abort

Step 2.

- provide fairness
- the first-to-commit thread commits consecutive txns

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synchronize threads

Evaluation

- Environment:
 - 2 Intel Xeon E5405 quad-core processors (8 cores in total),
 - 4GiB RAM
- TM support:
 - TinySTM 1.0.3.
- Benchmarks
 - STAMP [1]

^[1]Minh, C., Chung, J., Kozyrakis, C., Olukotun, K.: STAMP: Stanford transactional applications for multi-processing. In: Workload Characterization, IISWC 2008.



Evaluation - STAMP



Det-Parallel vs. Det-Serial (8 threads) BSC Microsoft Research Centre

Speedup: Bayes, Vacation Slowdown: Genome, Intruder, Kmeans, SSCA2

Conclusions

- Determinism:
 - to find and solve bugs easier
 - to provide fault- and attack- tolerance easier
- Different implementations of a deterministic system
- Provide deterministic and parallel execution of transactions
- Additional synchronization
 - provides strong determinism
 - lowers the benefit of parallel execution of transactions



Future work

- Eager Det-Parallel
 - to ensure progress of the first-to-commit transaction
 - to reduce the number of synchronization barriers
- Comparison with DTHREADS ^[2]
- Various schedule algorithms
- Determinism and diversity

^[2] T. Liu, C. Curtsinger, and E. D. Berger. DTHREADS: efficient deterministic Multithreading. In Proceedings of the 23rd ACM Symposium on Operating Systems Principles (SOSP '11).





Thank you for your attention!

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